

Coalition-based Resource Reciprocation Strategies for P2P Multimedia Broadcasting

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Abstract—In this paper, we consider a peer-to-peer (P2P) network, where multimedia streams are broadcast by matched peers based on their resource reciprocation profiles. We propose a new framework where each peer creates a coalition of matched peers with which it can exchange resources in order to improve its utility. The utility is determined based on explicit consideration of the peer's multimedia attributes and the quality derived by the peers' reciprocation behavior. We adopt the proportional bargaining solution to negotiate the upload bandwidth among the matched peers. Proportional bargaining allows to determine each peer's optimal (in a Pareto optimal sense) upload rates in a coalition in terms of its utility impact. The impact of an incoming peer on the coalition value, which represents the collective utility achieved by the peers in a coalition, is assessed by explicitly investigating the coalition value improvement. Finally, our results show that the proposed coalition-based resource reciprocation can improve the resource allocation/scheduling algorithms deployed in existing P2P systems such as BitTorrent and CoolStreaming. We also discuss how the proposed resource reciprocation approach can be implemented in other multimedia broadcasting applications.

Index Terms—Multimedia peer-to-peer (P2P) broadcast, P2P network, peer matching, service level, proportional bargaining solution, coalitions.

I. INTRODUCTION

P2P applications (e.g., [1], [2]) have become increasingly popular and represent a large majority of the traffic currently transmitted over the Internet. One of the unique aspects of P2P networks stems from its flexible and distributed nature, where each peer can act as both server and client [3]. Due to these characteristics, it has been recently proposed to use P2P networks for multimedia streaming and broadcasting [4]–[7]. Moreover, several media streaming and broadcasting systems have been developed for P2P networks using different approaches such as tree-based or data-driven approaches (e.g., [8]–[10]). While the above mentioned multimedia streaming and broadcasting systems over P2P networks have been successfully developed, the focus of P2P for multimedia is on developing efficient streaming solutions given a certain P2P network.

In this paper, we consider a chunk-based and data-driven multimedia P2P broadcasting system such as CoolStreaming [4] and Chainsaw [5], which adopt pull-based techniques [5], [6]. In this system, a multimedia stream is divided into chunks, which are broadcast over the P2P network. The peers possess several chunks and these chunks are shared by

the interested¹ peers. While this approach has been successfully deployed in real-time multimedia broadcasting over P2P networks, key challenges such as incentives for cooperation and designing efficient resource-reciprocation solutions among the autonomous peers still remain open. Specifically, pull-based techniques assume that the peers in the P2P network are altruistic and they provide their available chunks whenever requested.

Several works such as [11], [12] propose how to design the incentives for cooperation in general P2P networks. For example, a general file download solution such as BitTorrent adopts a *tit-for-tat* (TFT) algorithm (or *choking* algorithm) [11]. This algorithm favors peers that have higher upload rates and penalizes free-riders, thereby encouraging the peers to allocate more upload bandwidth [13]. While this algorithm can provide incentives for cooperation, it does not consider the exact resource reciprocation. Rather, a peer equally allocates its available bandwidth to a fixed number of peers (i.e., leechers) having the highest upload rates [11]. This approach has been shown to be inefficient for multimedia applications [14]. Hence, this solution cannot be directly deployed for multimedia broadcasting as it does not efficiently consider important issues related to the quality of service (QoS) requirements such as the derived quality and delay sensitivity. In order to address these issues for multimedia broadcasting, several works such as [15], [16] have been proposed. While they successfully provide incentives for cooperation in multimedia broadcasting, the resource reciprocation among peers is not explicitly considered.

In this paper, to explicitly consider the resource reciprocation among the interested multimedia peers, we discriminate the peers based on their multimedia attributes such as the derived quality, which is neglected by existing solutions. For instance, the TFT strategy in BitTorrent allows a peer to *choke*² the already associated peers at any rechoke period if they find new peers that can provide higher upload rates. Thus, TFT can only consider provisional connections among the associated peers. Hence, this strategy for multimedia content sharing cannot guarantee an acceptable quality for broadcast at all times. To overcome this issue, we introduce the service level, which represents the level of guaranteed bandwidth that the peer promises to sustain to the matched peers in its coalition. We show that a higher service level can reduce the quality fluctuations in P2P based multimedia broadcasting.

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¹In [11], it is said that peer A is *interested* in peer B when B has chunks of the content that A would like to possess.

²It is said that peer A is *choked* by peer B when B decides not to send any data to A.

A peer can create a *coalition* with its matched peers, which agree to share their chunks. In the created coalitions, peers negotiate their upload bandwidth based on their contributions to the coalitions, while explicitly considering their quality impact. We define appropriate utility functions to consider the derived multimedia qualities and the negotiated download/upload rates. Hence, utility-based resource negotiation enables peers to determine how much upload bandwidth they should provide, while considering the resulting quality impact. In order to achieve this, we propose to deploy the proportional bargaining solution (PBS) [17].

This paper is organized as follows. In Section II, we propose a framework for multimedia broadcasting over the P2P network and abstract the attributes of peers based on their multimedia characteristics. In Section III, we discuss the cooperative behavior and the corresponding achievable utilities. In Section IV, we discuss the upload bandwidth reciprocation among the peers and its resulting utilities determined by PBS. In Section V, we study how much the peers benefit by making coalitions and analyze the impact of their interactions in terms of the coalition value. Simulation results are provided in Section VI and conclusions are drawn in Section VII.

II. COALITION FORMATION BASED ON MULTIMEDIA PEERS' ATTRIBUTES

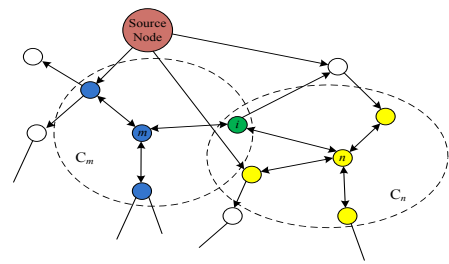
In this section, we discuss the attributes that need to be considered for multimedia applications and define the peer's utility function.

A. Coalition Formation

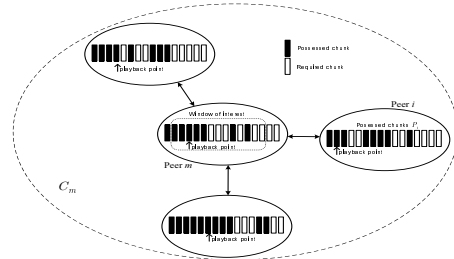
When a peer joins the P2P network, the first step for the incoming peer is to identify the peers that are interested in the same content. As discussed in [4], the incoming peer can contact a source node, which generates/broadcasts the new multimedia content, and the peer can be redirected to the other peers that are already associated with the source node in the overlay. Hence, the peers that are interested in the same content will interact with each other by sharing (i.e., downloading and uploading) chunks of the interested content. The detailed process to create and maintain the connections among the peers can be found in e.g., [4], [5]. In this paper, we consider the case where a source node distributes multimedia content and the peers that are interested in it interact with each other by sharing their content and negotiating their resources. This can be easily extended to the case when multiple multimedia contents from multiple source nodes are simultaneously broadcast.

Our focus is on the resource reciprocation among the peers that are interested in each other³. When two peers are interested in each other, they are referred to as *matched* peers. More specifically, any two peers m and n are matched if peer m is interested in the chunks that peer n possesses

³If a peer m is interested in n , while n is not interested m , then we assume that peer n provides the minimum upload bandwidth, which is similar to the optimistic unchoke in BitTorrent [11]. This operation can help to maintain the coalitions and to increase the P2P network capacity as peer n can contribute its chunks in the future interactions.



(a) An illustration of coalitions (C_m and C_n)



(b) Peers matched to peer m in C_m and their chunks

Fig. 1. Illustrative examples of coalitions and chunks.

and peer n is also interested in the chunks that peer m possesses. Peer m includes in its coalition C_m its matched peers with which it will exchange chunks. In the process of chunk exchanges, resource reciprocation is necessary for the peers to improve their utilities based on their limited upload bandwidths (i.e., resources). An illustrative example of coalitions in a P2P network is shown in Fig. 1 (a). In this example, two coalitions C_m and C_n , formed by peer m and peer n with their matched peers, are depicted. Moreover, the chunks possessed and required by the peers in coalition C_m are illustrated in Fig. 1 (b). A peer maintains a window of interest (or an active buffer) as in [4], [5], which represents the range of chunks that the peer is interested in obtaining at the current time. As Fig. 1 (b) shows, the chunks required for peer m are available in its coalition C_m .

B. Models for Multimedia Peer's Attributes

The multimedia peers in the P2P network can be characterized by their possessed and demanded chunks, their preferences for the content, available bandwidth, as well as the cost for upload [18]–[21]. Moreover, a quality function is required to assess the achievable multimedia quality based on the download rates, and a minimum download rate for a demanded multimedia content is required to successfully decode the content [22]. We define a peer's service level as its commitment to sustaining an upload bandwidth for its matched peers. In the considered P2P-based multimedia broadcasting, we can abstract the attributes of peer i as

$$\mathbf{A}_i = \{D_i, \rho_i, R_{D_i}^{req}, P_i, L_i, Q_i(\cdot), \gamma_i, c_i\}, \quad (1)$$

where D_i denotes a multimedia content in which peer i is interested with its preference ρ_i and $R_{D_i}^{req}$ is the minimum required rate to successfully decode D_i . If this condition is not satisfied, a peer cannot decode the downloaded multimedia

chunks, and hence, it achieves no utility. P_i denotes the available chunks for peer i , which will be continuously exchanged among the associated peers. This is similar to the Buffer Map deployed in [4]. $L_i (\leq L_i^{MAX})$ denotes the available upload bandwidth⁴ of peer i . Note that the available upload bandwidth of a peer will decrease as it associates with more peers. The achieved quality given the download rate R_D is denoted by $Q_i(R_D)$, which can be represented by the Peak Signal to Noise Ratio (PSNR). γ_i is the service level, which will be discussed later, and c_i denotes the per unit cost of uploading rate that can be determined as in [21].

We define a utility function that captures this behavior of multimedia peers. The utility function of peer i , which is downloading content D_i at rate R_{D_i} and providing upload rate R_{U_i} to other peers, is defined as:

$$U_i(R_{D_i}, R_{U_i}) = \begin{cases} 0, & \text{if } R_{D_i} < R_{D_i}^{req}, \\ \rho_i Q_i(R_{D_i}) - c_i R_{U_i}, & \text{otherwise,} \end{cases} \quad (2)$$

where a non-negative constant $R_{D_i}^{req}$ represents the specific minimum required rates to decode the video sequence. Peer i 's utility increases as it derives a higher quality by downloading at higher rates. The utility also increases as peer i reduces its upload bandwidth. Download rate R_{D_i} is a function of R_{U_i} , i.e., $R_{D_i}(R_{U_i})$, since they are determined based on the resource reciprocation among the associated peers⁵. The resource reciprocation is investigated in Section III. Note that different utilities can be defined and employed in our framework depending on the P2P applications.

The service level⁶ γ_i ($0 \leq \gamma_i \leq 1$) represents peer i 's level of commitment to maintain the upload bandwidth promised to the associated peers. For instance, the available upload bandwidth of peer i matched to the other peers is determined by

$$L_i = L_i^{MAX} - \gamma_i \cdot \sum_{l \in C_i \setminus \{i\}} R_{U_l}, \quad (3)$$

where γ_i is the service level of peer i , and R_{U_l} denotes peer i 's current upload bandwidth for peer $l \in C_i \setminus \{i\}$. This expression shows that a higher value of service level of a peer will result in a lower available upload bandwidth for the future resource negotiation with incoming peers, as it sustains an upload bandwidth for the already established connections. For instance, peer i with $\gamma_i = 0$ uses its maximum available upload bandwidth whenever new incoming peers join C_i , while it only uses the remaining upload bandwidth if $\gamma_i = 1$. The effect of the service level on making coalitions and negotiating upload bandwidths will be discussed in Section V. Importantly, it will be shown that a higher service level provides more stable multimedia quality.

Note that other characteristics of P2P networks such as reliability, contract, and reputation can also be included as

attributes of peers. However, we do not consider them in this paper, as they are already discussed in the prior literature [23], [25], and they can be incorporated in our framework.

III. CHARACTERIZATION OF PEERS' INTERACTIONS

In this section, we consider the utility-based interactions between autonomous and self-interested peers in a P2P network.

A. Motivating Example: Incentives for Peers' Collaboration

As motivating examples, we discuss simple cases where two peers do not cooperate or overly cooperate. Note that this derivation can be easily extended to the general N -peer case. In this analysis, we consider two matched peers m and n .

1) *Non-cooperative Interactions*: Based on the definition of utility function as in (2), the utilities of peer m and n can be expressed as

$$\begin{aligned} U_m(R_{U_n}, R_{U_m}) &= \rho_m Q_m(R_{U_n}) - c_m R_{U_m}, \\ U_n(R_{U_m}, R_{U_n}) &= \rho_n Q_n(R_{U_m}) - c_n R_{U_n}. \end{aligned}$$

If both peers are non-cooperative, each peer only tries to improve its utility without considering the impact on the other peers' utilities. Hence, peer m can improve its utility by reducing its upload rate by ΔR_{U_m} , and peer n can also reduce its upload rate by ΔR_{U_n} as a response of peer m 's upload rate change. The resulting utilities of two peers become $U_m(R_{U_n} - \Delta R_{U_n}, R_{U_m} - \Delta R_{U_m})$ and $U_n(R_{U_m} - \Delta R_{U_m}, R_{U_n} - \Delta R_{U_n})$. To minimize the penalties due to the cost for upload, two peers will continuously reduce their upload rate, which leads to the bandwidth reciprocation $\Delta R_{U_m} = R_{U_m}$ and $\Delta R_{U_n} = R_{U_n}$, and thus, zero utilities (i.e., $U_m = 0$, $U_n = 0$). Thus, the non-cooperative behavior of peers in a P2P network eventually results in zero utilities for every peer. It can be interpreted that if there is no cooperation among peers, then the collapse of a P2P network is possible [26], i.e., no utility can be provided to peers through a P2P network. A well-known example for this non-cooperative behavior of peers is the free-riding [26] in P2P networks.

2) *Overly Cooperative Interactions*: Alternatively, we consider the case where peers are completely cooperative, by providing their entire upload bandwidth to other peer's download. Let L_m^{MAX} and L_n^{MAX} be the maximum available upload bandwidth for peers m and n , respectively. Since they completely collaborate by providing their maximum available upload bandwidth, their achievable utilities are expressed as $U_m(L_n^{MAX}, L_m^{MAX})$ and $U_n(L_m^{MAX}, L_n^{MAX})$, respectively. Since these utilities can be negative depending on their costs, quality functions, and the provided maximum upload bandwidth, the complete cooperation among peers does not always benefit them. Rather, it may penalize the participating peers in terms of the derived utility. Moreover, this type of over altruistic cooperation can also prevent incoming peers which join the P2P network at a later time from associating with these peers, as the available upload bandwidth of the peers in the network is already exhausted. This may lead to inefficient bandwidth utilization for peers as well as the P2P system in terms of the derived utilities.

⁴In this paper, we focus on the upload bandwidth allocation since the maximum upload bandwidth L_i^{MAX} of peer i is typically smaller than the maximum download bandwidth, e.g., ADSL, cable network, etc.

⁵In this paper, we use notation R_{D_i} instead of $R_{D_i}(R_{U_i})$ for simplicity.

⁶The service level can be determined based on peers' reputation [23], experience, or characteristics such as risk-averse or risk-lover [24]. While we do not explicitly examine algorithms to determine this parameter dynamically, the impact on coalition formation and upload bandwidth negotiation can be investigated by considering two extreme cases.

Therefore, the cooperative negotiation for determining the upload bandwidth among peers is essential for improving both each peer's utility and the total utility in P2P network.

B. Example of Collaboration between Two Peers

In this section, we analyze the collaborative two-peer's interaction and its achievable utility. Unlike the non-cooperative types of peers, the matched peers exchange information that is required to establish their cooperation. In this simple cooperative two-peer scenario, we assume that two peers m and n agree before making their coalition to provide their upload bandwidth as long as the allocated bandwidth enables them to achieve positive utilities. Note that each peer achieves no utility (i.e., zero utility) if they are not associated with each other. Based on this agreement, peer m can first provide its upload bandwidth $R_{U_m} (\geq R_{D_n}^{req})$ to peer n [27], which results in

$$U_m = -c_m R_{U_m}, \quad U_n = \rho_n Q_n(R_{U_m}).$$

After receiving R_{U_m} from peer m , peer n can now provide its upload bandwidth R_{U_n} if it does not make its achievable utility negative (if it is not possible, they cannot interact cooperatively). Hence, $R_{U_n} (\geq R_{D_m}^{req})$ needs to satisfy

$$U_n = \rho_n Q_n(R_{U_m}) - c_n R_{U_n} \geq 0,$$

and it leads to the maximum upload rate that peer n can provide, which is

$$R_{U_n} \leq \rho_n Q_n(R_{U_m}) / c_n. \quad (4)$$

Moreover, the upload rate from peer n should also guarantee that peer m can derive a positive utility. Hence,

$$U_m = \rho_m Q_m(R_{U_n}) - c_m R_{U_m} \geq 0,$$

which leads to the minimum upload rate,

$$R_{U_n} \geq Q_m^{-1}(c_m R_{U_m} / \rho_m). \quad (5)$$

From (4), (5), and the required minimum rates, peer n can provide its upload rate R_{U_n} in the range of

$$\max \left\{ R_{D_m}^{req}, Q_m^{-1} \left(\frac{c_m R_{U_m}}{\rho_m} \right) \right\} \leq R_{U_n} \leq \frac{\rho_n Q_n(R_{U_m})}{c_n}. \quad (6)$$

Summarizing, (6) implies that the upload rate of peer n is upper-bounded by the rate that achieves the minimum utility of peer n and lower-bounded by the rate that achieves the minimum utility of peer m , which are a function of the upload rate of peer m . Note that R_{U_n} in (6) can be parameterized by the variable θ_n ($0 \leq \theta_n \leq 1$):

$$R_{U_n}(R_{U_m}) = \theta_n \cdot \left[\max \left\{ R_{D_m}^{req}, Q_m^{-1} \left(\frac{c_m R_{U_m}}{\rho_m} \right) \right\} \right] + (1 - \theta_n) \cdot \left[\frac{\rho_n Q_n(R_{U_m})}{c_n} \right], \quad (7)$$

where $R_{U_m} \geq R_{D_n}^{req}$. As Fig. 2 illustrates, it is easily observed that there is a conflict of interest between two peers' utilities according to their upload rates. For instance, if $\theta_n = 1$, where peer n supports the minimum upload rate to peer m , peer m can achieve the minimum utility 0 while peer n can

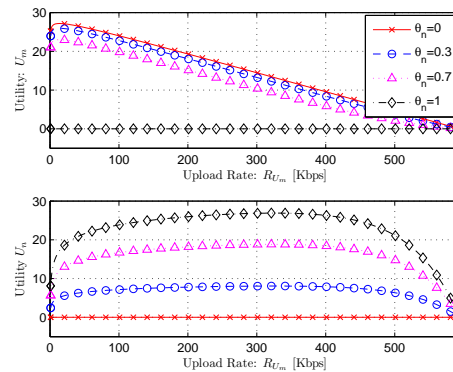


Fig. 2. Achievable utilities U_m and U_n for two peers with respect to peer m 's upload rate R_{U_m} . $\theta_n = 0$ induces the highest utility for peer m while peer n achieves the lowest utility, and vice versa. The shapes of achievable utilities are determined by their utility functions and peer-dependent cost.

achieve the highest achievable utility, and vice versa. Hence, variable θ_n can be viewed as an altruism parameter [28] of peer n , where $\theta_n = 1$ represents that peer n is the least altruistic for providing its upload bandwidth, while $\theta_n = 0$ represents that peer n is ideally altruistic. However, we note that in this paper, this variable is not a pre-determined peer-specific constant. Rather, it will be determined by the upload bandwidth negotiation, which will be discussed in Section IV. Hence, each peer's utility can be expressed as a function of R_{U_m} using (7): $U_m(R_{U_m}) = \rho_m Q_m(R_{U_n}(R_{U_m})) - c_m R_{U_m}$ and $U_n(R_{U_m}) = \rho_n Q_n(R_{U_m}) - c_n R_{U_n}(R_{U_m})$.

Based on the above analysis, we can summarize that, through cooperation, the peers can derive positive utilities. However, the utilities that each peer can derive are jointly connected as a function of each peer's upload rates and θ . Hence, it is imperative to determine the optimal upload rates of each peer in the coalition (e.g., determine R_{U_i} and θ_i ($i \in \{m, n\}$) in this example), while explicitly considering the derived utility and each peer's contribution to the coalition. This will be discussed in Section IV.

C. Cooperative Peers' Interactions

In this section, we generalize the interactions of peers sharing multimedia content in a P2P network. Since a peer interacts with several matched peers in its coalition, we focus on a one-to-many peers' interaction.

Let $C_1 = \{1, \dots, N\}$ be the coalition of peer 1, where peer 1 interacts with its $(N-1)$ matched peers while negotiating resource reciprocation. The utilities of the peers in the coalition are expressed as

$$U_1 = \rho_1 Q_1 \left(\sum_{i \in C_1 \setminus \{1\}} R_{U_{i1}} \right) - c_1 \left(\sum_{i \in C_1 \setminus \{1\}} R_{U_{i1}} \right), \quad (8)$$

$$U_i = \rho_i Q_i \left(R_{U_{i1}} + \sum_{l \in C_i \setminus \{1, i\}} R_{U_{li}} \right) - c_i \left(R_{U_{i1}} + \gamma_i \sum_{l \in C_i \setminus \{1, i\}} R_{U_{li}} \right), \quad \text{for all } i \in C_1 \setminus \{1\}, \quad (9)$$

where $R_{U_{ji}} (\geq R_{D_i}^{req})$ denotes the upload rate of peer j to peer i . Note that peer i in C_1 assumes that the resource

reciprocations of the associated peers in its coalition C_i are stationary as shown in (9). The expression of $B \setminus A$ for two sets represents the relative complement of A in B . Since peer i is associated with peer 1 as well as other peers in C_i , the achieved utility of peer i depends on the upload rates $R_{U_{1i}}$ provided by peer 1 and $R_{U_{li}}$ provided by the other peers $l \in C_i \setminus \{1, i\}$. It also depends on the provided upload rates from peer i to the associated peers. Note that the upload rates $R_{U_{i1}}$ that peer i can provide to peer 1 depends on the service level γ_i of peer i . If $\gamma_i = 0$, which implies that peers do not promise to sustain the already established connections in C_i , then peer i can provide its maximum upload bandwidth to peer 1, i.e., $0 \leq R_{U_{i1}} \leq L_i^{MAX}$. However, if $\gamma_i = 1$, peer i can provide only the remaining upload bandwidth to peer 1, i.e., $0 \leq R_{U_{i1}} \leq L_i$, while sustaining the already established connection in its coalition C_i .

Since every peer needs to achieve at least a minimum utility (i.e., zero utility), (8) and (9) can be rewritten as

$$R_{U_{i1}} \geq \max\{R_{D_1}^{req}, Q_1^{-1}\left(\frac{c_1}{\rho_1} \sum_{l \in C_1 \setminus \{1\}} R_{U_{1l}}\right) - \sum_{l \in C_1 \setminus \{1, i\}} R_{U_{1l}},\right\} \quad (10)$$

$$R_{U_{i1}} \leq \frac{\rho_i}{c_i} Q_i(R_{U_{1i}} + \sum_{l \in C_i \setminus \{1, i\}} R_{U_{li}}) - \sum_{l \in C_i \setminus \{1, i\}} R_{U_{il}} \cdot \gamma_i \triangleq R_{U_{i1}}^{MAX}, \text{ for all } i \in C_1 \setminus \{1\}. \quad (11)$$

Eq. (10) can be expressed using $R_{U_{i1}}^{MAX}$ in (11) as

$$R_{U_{i1}} \geq \max\{R_{D_1}^{req}, Q_1^{-1}\left(\frac{c_1(\sum_{l \in C_1 \setminus \{1\}} R_{U_{1l}})}{\rho_1}\right) - \sum_{l \in C_1 \setminus \{1, i\}} R_{U_{1l}}^{MAX} \triangleq R_{U_{i1}}^{min}\right\} \quad (12)$$

Therefore, the upload rate $R_{U_{i1}}$ can be expressed as

$$R_{U_{i1}} = \theta_i \cdot R_{U_{i1}}^{min} + (1 - \theta_i) \cdot R_{U_{i1}}^{MAX}, \quad (13)$$

with variable θ_i ($0 \leq \theta_i \leq 1$). Note that $R_{U_{i1}}^{MAX}$ is a function of upload rates $R_{U_{1i}}$ from peer 1 to peer i and parameters of peer i in C_i (i.e., upload/download rates in C_i and γ_i). $R_{U_{i1}}^{min}$ is also a function of upload rates $R_{U_{1l}}$, $l \in C_1 \setminus \{1\}$ from peer 1 to its peers in C_1 , and coalition parameters of peer l , $l \in C_1 \setminus \{1\}$. Therefore, by substituting (13) into (8) and (9), the achievable utilities in C_1 can be expressed as a function of peer 1's upload rates to its coalition peers given the other coalition parameters.

Hence, it can be concluded that if the feasible upload rate pairs satisfying (11) and (12) can be found, the peers' cooperative behavior finally benefits the participating peers in this cooperative interaction. Therefore, it is essential for peers to form coalitions that enable peers to interact with each other cooperatively. The remaining question is hence how to negotiate an agreement for determining the upload rates of all peers after making a coalition (i.e., determining θ_l for all $l \in C$ in a coalition C). In the following sections, we resolve this problem by relying on axiomatic bargaining theory among the peers. Note that the proposed bargaining solution provides

axioms that can be dictated by the adopted P2P protocol for the resource reciprocation in the coalitions. Hence, there is no need for an iterative bargaining process among peers.

IV. COALITION BASED UPLOAD BANDWIDTH NEGOTIATION

In this section, we discuss the upload bandwidth negotiation and determine the optimal upload rates of the peers in a coalition based on the derived utility and each peer's contribution to the coalition.

A. Resource Reciprocation and Resulting Utilities

In order to determine the optimal upload rates of peers based on the derived utility and each peer's contribution to the coalition, it is required to identify the feasible utility set. For this analysis, we again assume that peer 1 interacts with the other $(N - 1)$ peers in C_1 . The feasible utility set $\mathbf{S} \subset \mathbb{R}^N$ can be obtained given the utility functions of peers and their interactions for the upload rates, which is expressed as

$$\mathbf{S} = \{(U_1(\sum_{l=2}^N R_{U_{1l}}, \sum_{l=2}^N R_{U_{1l}}), \mathbf{U}) \mid R_{U_{1l}} \in \mathbf{R}_{U_{1l}}, R_{U_{1l}} \in \mathbf{R}_{U_{1l}}, \forall l\}, \quad (14)$$

where $\mathbf{R}_{U_{1l}} = \{R_{U_{1l}} \mid 0 \leq R_{U_{1l}} \leq L_1, \sum_{l \in C_1} R_{U_{1l}} \leq L_1\}$, $\mathbf{R}_{U_{1l}} = \{R_{U_{1l}} \mid 0 \leq R_{U_{1l}} \leq L_l, R_{U_{1l}} + \gamma_l \sum_{h \in C_l} R_{U_{lh}} \leq L_l, 0 \leq \gamma_l \leq 1\}$ for all l ($2 \leq l \leq N$), and $\mathbf{U} = [u_2, \dots, u_N]$, where $u_h = U_h(R_{U_{1h}} + \sum_{l \in C_h \setminus \{h\}} R_{U_{lh}}, R_{U_{h1}} + \sum_{l \in C_h \setminus \{h\}} R_{U_{hl}})$ for $2 \leq h \leq N$. The feasible utility set \mathbf{S} includes all the possible interactions of peers, i.e., cooperative interactions as well as non-cooperative interactions. Moreover, each peer's minimum required utility⁷ $\mathbf{d} \subset \mathbb{R}^N$ can also be identified and considered in the upload bandwidth negotiation. The outcome of the upload bandwidth negotiation needs to be Pareto optimal and yield higher utilities than the minimum required quality. Thus, peers' agreement on their derived utilities is located in set⁸ $\mathbf{B} \subseteq \mathbf{S}$, expressed as

$$\mathbf{B} = \{\partial \mathbf{S} \cap \{\mathbf{U} \mid \mathbf{U} \geq \mathbf{d} \text{ for all } \mathbf{U} \in \mathbb{R}^N\}\}, \quad (15)$$

where $\partial \mathbf{S}$ denotes the boundary [17] of feasible utility set \mathbf{S} . An illustrative example of utility pairs derived by the results of two-peer's cooperative interactions is shown in Fig. 3. Different upload rates result in different utility pairs and their maximal values are elements of the bargaining set.

B. Resource Reciprocation based on Contributions

We now discuss how the peers can agree on a unique utility point. Given feasible utility set \mathbf{S} in (14) and the minimum required utility \mathbf{d} , peers can agree on a unique utility point based on a vector $\mathbf{p} = (p_1, \dots, p_N)$, referred to as *inter-peer utility comparison ratio vector*. The vector \mathbf{p} represents each peer's contribution to their coalition and can be adjusted based on the goal of the peers in the coalition. For example, if the available upload bandwidth of peers are considered as

⁷This is called *disagreement point* in axiomatic bargaining theory [24].

⁸This set is called *bargaining set* [24].

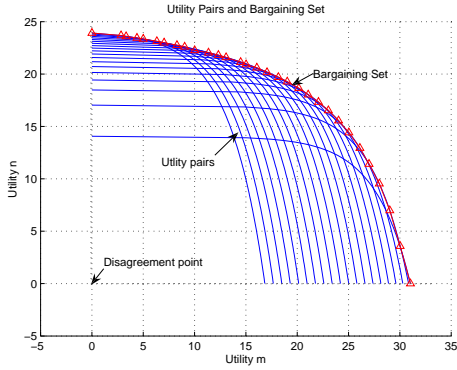


Fig. 3. Illustrative examples for utility pairs derived by cooperative interactions of peers.

their contributions, vector \mathbf{p} can be determined as $p_i = L_i$, $i = 1, \dots, N$. Alternatively, the number of chunks that are requested by the matched peers in a coalition can also be considered as contributions, which can determine \mathbf{p} .

By deploying proportional bargaining solution (PBS) [17] to the bargaining problem (\mathbf{S}, \mathbf{d}) with the vector \mathbf{p} , a unique utility point $\mathbf{U}^* \in \mathbf{B}$ can be thereby determined as

$$F(\mathbf{S}, \mathbf{d}) = \mathbf{U}^* = (U_1^*, \dots, U_N^*), \quad (16)$$

where F denotes the PBS and $U_i^* = [F(\mathbf{S}, \mathbf{d})]_i$. Note that the determined utility point \mathbf{U}^* based on the PBS F satisfies

$$U_1^*/p_1 = \dots = U_N^*/p_N. \quad (17)$$

Hence, vector \mathbf{p} can be interpreted as the weight of the peers based on their contributions to the coalition. Several properties of PBS can be found in [17]. Among them, *monotonicity*⁹ will be used to analyze the coalition value dynamics as peers join the coalitions (see Section V). Note that given vector \mathbf{p} , the unique utility pair selected by PBS is optimal (in the sense of Pareto optimality).

Determining a unique solution \mathbf{U}^* by PBS can be efficiently performed using numerical tools such as the bisection method without identifying the entire feasible utility set [14]. Although computing PBS may incur higher initial delay for resource reciprocation process than heuristic approaches deployed in [4], random packet selection approach [5], or the TFT strategy, the proposed approach can efficiently utilize multiple peers' resources. Hence, the proposed approach can improve the video quality and provide a steady video quality, which is desired for multimedia broadcasting. Hence, overall delay for multimedia broadcasting can be reduced. These tradeoffs are illustrated in Section VI.

Algorithm 1 summarizes the negotiation of the upload bandwidth based on PBS in a coalition.

V. COALITION VALUE AND ITS DYNAMICS

In this section, we study how much the peers benefit by making coalitions and analyze the impact of their interactions in terms of the coalition value.

⁹*Monotonicity*: For any two feasible utility sets \mathbf{S}' and \mathbf{S} , if $\mathbf{S}' \supset \mathbf{S}$ then $[F(\mathbf{S}, \mathbf{d})]_i \geq [F(\mathbf{S}', \mathbf{d})]_i$ for all i .

Algorithm 1 Upload Bandwidth Negotiation in C_i

- 1: *Form Feasible Utility Set \mathbf{S}* : From peers' attributes \mathbf{A}_l , $l \in C_i$, peers form \mathbf{S} and the disagreement point \mathbf{d} .
- 2: *Determine inter-peer utility comparison ratio vector \mathbf{p}* : Based on the peers' contribution to the coalition, vector \mathbf{p} for PBS is determined.
- 3: *Determine Operating Utility Point \mathbf{U}^** : A unique utility point \mathbf{U}^* is determined based on the PBS with \mathbf{p} ; $\mathbf{U}^* = F(\mathbf{S}, \mathbf{d})$.
- 4: *Determine Upload Bandwidth*: Given \mathbf{U}^* , peers determine their upload bandwidth; $R_{U_i}^* = U_i^{-1}(U_i^*)$ for all l .

A. Service Level and Coalition Value Dynamics

Coalition value is the total utility derived by the peers in a coalition and represents the benefit of cooperation among the peers in the coalition. Let $C_i^{(k)}$ be the coalition of peer i at round¹⁰ k . Initially, peer i exists as a self-coalition at round 0, i.e., $C_i^{(0)} = \{i\}$ with $\mathbf{d}_i^{(0)} = 0$. Incoming peers matched to peer i can join C_i . In $C_i^{(k)}$, the coalition value for $C_i^{(k)}$ determined by PBS is given by

$$v(C_i^{(k)}) = \sum_{l \in C_i^{(k)}} [F(\mathbf{S}_i^{(k)}, \mathbf{d}_i^{(k)})]_l,$$

where $\mathbf{S}_i^{(k)}$ is the feasible utility set formed by the interaction among the peers in $C_i^{(k)}$ and $\mathbf{d}_i^{(k)}$ is their disagreement point.

Suppose that an incoming peer j matched to peer i joins C_i at round $k+1$, i.e., $C_i^{(k+1)} = C_i^{(k)} \cup \{j\}$. For the PBS with peer j at round $(k+1)$, the disagreement point of all peers $l \in C_i^{(k)}$ is updated by $d_l^{(k+1)} \in \mathbf{d}_i^{(k+1)}$, expressed as

$$d_l^{(k+1)} = d_l^{(k)} + U_l \left(\gamma_i \cdot U_i^{-1} \left([F(\mathbf{S}_i^{(k)} - \mathbf{d}_i^{(k)}, \mathbf{0})]_i \right) \right) \quad (18)$$

for $l \in C_i^{(k)} \setminus \{i\}$, and

$$d_i^{(k+1)} = U_i \left(\sum_{l \in C_i^{(k)} \setminus \{i\}} \gamma_l \cdot U_l^{-1} ([F(\mathbf{S}_i^{(k)} - \mathbf{d}_i^{(k)}, \mathbf{0})]_l) \right). \quad (19)$$

where γ_l is the service level of peer l in C_i . Moreover, $d_j^{(k+1)} = 0$, since peer j is an incoming peer.

As shown in (3), (18), and (19), the service levels significantly impact the resource reciprocation among the peers in their coalitions. For example, if peers in a coalition have $\gamma = 0$, they try to find better matched peers that can provide higher upload rates at every round as they are not obligated to sustain any upload bandwidth for currently associated peers. While this can lead to a higher achievable utility for peers depending on the associated peers, they can suffer from frequent connection loss, which can lead to significant fluctuations of the derived utility. Moreover, the complexity required for PBS computation can increase exponentially since the number of peers that can be associated with a peer can increase significantly. Hence, the interaction of peers with $\gamma = 0$ do not guarantee a steady level of quality for multimedia applications. However, if peers have $\gamma = 1$, they sustain the established

¹⁰The term *round* is adopted from [11] to represent the time when PBS is executed. In this paper, rounds are equivalent to the instances of peers' joining or leaving coalitions.

connections unless the associated peers leave the P2P network. Peers in a coalition can derive less achievable utility, but they are able to attain a stable achievable utility, leading to quality guarantee for multimedia. This is an important feature for multimedia content distribution over P2P networks. Hence, the service level needs to be determined considering these tradeoffs. In the rest of the analysis, we assume that the service level of each peer is $\gamma = 1$.

B. Peers' Joining and Coalition Value Dynamics

In order to quantify the impact of peer j 's joining $C_i^{(k)}$ on the coalition value, the coalition value improvement needs to be investigated. Since $\gamma_l = 1$ for all $l \in C_i^{(k)}$, we have

$$\begin{aligned} v(C_i^{(k+1)}) &= v(C_i^{(k)} \cup \{j\}) = \sum_{l \in C_i^{(k+1)}} \left[F(\mathbf{s}_i^{(k+1)}, \mathbf{d}_i^{(k+1)}) \right]_l \\ &= \sum_{l \in C_i^{(k+1)}} d_l^{(k+1)} + \sum_{l \in C_i^{(k+1)}} \left[F(\mathbf{s}_i^{(k+1)} - \mathbf{d}_i^{(k+1)}, \mathbf{0}) \right]_l. \end{aligned}$$

Similarly,

$$v(C_i^{(k)}) = \sum_{l \in C_i^{(k)}} d_l^{(k)} + \sum_{l \in C_i^{(k)}} \left[F(\mathbf{s}_i^{(k)} - \mathbf{d}_i^{(k)}, \mathbf{0}) \right]_l.$$

Hence, the improvement of coalition value $v(C_i^{(k)})$ by an additional matched peer j can be expressed as

$$\Delta v_j(C_i^{(k)}) = \sum_l \left[F(\mathbf{s}_i^{(k+1)} - \mathbf{d}_i^{(k+1)}, \mathbf{0}) \right]_l. \quad (20)$$

Since $\sum_l [F(\mathbf{s}_i^{(k+1)} - \mathbf{d}_i^{(k+1)}, \mathbf{0})]_l \geq 0$, the joining of a matched peer does not decrease the coalition value.

We also show that the coalition value improvement is a non-increasing function as the number of peers with similar attributes increases. Let peer i be in a coalition at round $k-1$, and peer j_k and peer j_{k+1} join the coalition at round k and $k+1$, respectively. Since peer i uses the residual bandwidth that is left after bargaining at round k with peer j_k as upload rate for peer j_{k+1} at round $k+1$, and the disagreement point of peer i increases (i.e., $\mathbf{d}_i^{(k)} \leq \mathbf{d}_i^{(k+1)}$), we have

$$\mathbf{S}_i^{(k)} - \mathbf{d}_i^{(k)} \supseteq \mathbf{S}_i^{(k+1)} - \mathbf{d}_i^{(k+1)}, \quad (21)$$

where $\mathbf{S}_i^{(k)}$ and $\mathbf{S}_i^{(k+1)}$ are the feasible utility sets formed at round k and $k+1$, respectively. By the monotonicity of the PBS,

$$F(\mathbf{S}_i^{(k)} - \mathbf{d}_i^{(k)}, \mathbf{0}) \geq F(\mathbf{S}_i^{(k+1)} - \mathbf{d}_i^{(k+1)}, \mathbf{0}), \quad (22)$$

and therefore,

$$\Delta v_{j_k}(C_i^{(k)}) \geq \Delta v_{j_{k+1}}(C_i^{(k+1)}). \quad (23)$$

Thus, the coalition value improvement is a non-increasing function. Since the maximum upload bandwidth is generally finite, we have $\Delta v_{j_k}(C_i^{(k)}) > \Delta v_{j_{k+1}}(C_i^{(k+1)})$, which implies that the coalition value improvement by incoming peer j_k consequently converges to zero as $k \rightarrow \infty$, i.e., $\lim_{k \rightarrow \infty} \Delta v_{j_k}(C_i^{(k)}) = 0$. This upload bandwidth negotiation can be easily applied iteratively at each round, as multiple peers join the coalition.

VI. SIMULATION RESULTS

In this section, we provide simulation results of the proposed framework for matching peers for multimedia broadcasting over P2P networks. In order to illustrate the potential impact of the proposed coalition-based resource reciprocation on multimedia P2P broadcasting networks, several simulations are performed based on the PlanetLab experimental platform as in [11]. We assume that the broadcast video files are at CIF (352×288) resolution, 30 frames/sec, as well as encoded in a prioritized manner using the H.264/AVC encoder [22]. The encoded video files are partitioned into chunks that have uniform size of 50 Kbits. Peers download first the chunks that have the highest impact on the video quality. In these experiments, a single video file has 100 seconds duration, which was obtained by concatenating 10 identical MPEG test sequences. We assume that peers can change their available upload bandwidth, or disconnect at each round. In our simulations, 1000 peers are registered in the P2P network with various multimedia attributes. Relevant peers' attributes are presented in each simulation.

A. Effect of Service Level

We investigate two cases of service levels, $\gamma_i = 1$ and $\gamma_i = 0$ for all i in a coalition. Recall that if $\gamma_i = 1$ for all peers in a coalition, they maintain their established connections and negotiate their remaining bandwidth with incoming peers. If $\gamma_i = 0$ for all peers in a coalition, peers do not promise to sustain their established connection and renegotiate their available upload bandwidth at all times. In order to investigate the impact of the service level, we assume that $p_i = p_j$ for all $i \neq j$ in vector \mathbf{p} . We focus on coalition 1 (C_1) which is created by peer 1 and its matched incoming peers that join the coalition sequentially. Two separate simulations are performed with two different video sequences: *Foreman* and *Coastguard*. The maximum upload bandwidth of each peer is given by L_i^{MAX} [kbps] ($1 \leq i \leq 8$) = {350, 400, 300, 350, 350, 600, 280, 350}. Simulation results are shown in Fig. 4 and Table I.

Fig 4 shows the achieved quality of peer 1 when it downloads the *Foreman* sequence or *Coastguard* video sequences. As discussed above, the achieved quality with $\gamma = 1$ provides an stable quality over time because the associated peers sustain their upload bandwidth. However, while the achieved quality with $\gamma = 0$ can be higher than that with $\gamma = 1$ (e.g., first 12 seconds), significant fluctuation in the achieved quality can occur, which is not desirable for multimedia broadcasting (e.g., approximately a 10 dB and 6 dB difference in terms of PSNR, respectively). Note that a larger fluctuation in PSNR generally results in a higher fluctuation in perceptual quality [29], thereby leading to significant perceptual quality degradation [30]. Hence, $\gamma = 1$ is preferred for the considered multimedia P2P broadcasting system.

B. Comparison with Other Resource Reciprocation Strategies

In this section, we illustrate the achievable utilities of peers in a coalition based on different resource reciprocation strategies. We consider several different P2P network scenarios: (i)

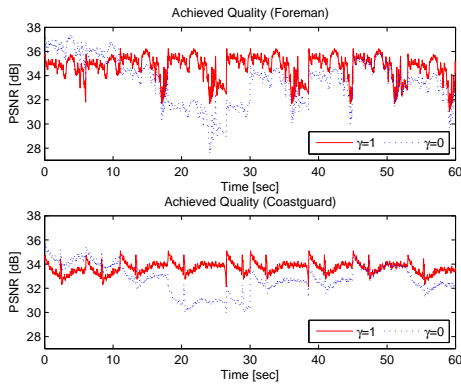


Fig. 4. Achieved quality with different service levels in a coalition.

TABLE I

ACHIEVED QUALITIES WITH DIFFERENT SERVICE LEVELS IN COALITION

| Foreman | Avg. PSNR | MAX (PSNR) | min (PSNR) | Difference |
|--------------|-----------|------------|------------|------------|
| $\gamma = 0$ | 33.5 dB | 37.4 dB | 27.5 dB | 9.9 dB |
| $\gamma = 1$ | 34.7 dB | 36.2 dB | 31.0 dB | 5.2 dB |
| Coastguard | Avg. PSNR | MAX (PSNR) | min (PSNR) | Difference |
| $\gamma = 0$ | 32.7 dB | 35.8 dB | 29.9 dB | 5.9 dB |
| $\gamma = 1$ | 33.6 dB | 35.0 dB | 32.1 dB | 2.9 dB |

TABLE II

ACHIEVABLE INDIVIDUAL QUALITIES IN DIFFERENT SCENARIOS

| PSNR [dB] | Equal Bandwidth Division | | PBS with \mathbf{p}_E |
|------------|---------------------------------|---------------|-------------------------|
| | Scenario (i) | Scenario (ii) | Scenario (iii) |
| Peer 1 | 38.37 | 38.37 | 39.81 ($U_1=35.00$) |
| Peer 2 | 32.19 | 33.95 | 36.85 ($U_2=35.00$) |
| Peer 3 | 30.99 | 32.57 | 35.18 ($U_3=35.00$) |
| Free Rider | 32.19 | - | - |
| PSNR [dB] | Proportional Bandwidth Division | | PBS with \mathbf{p}_P |
| | Scenario (i) | Scenario (ii) | Scenario (iii) |
| Peer 1 | 38.37 | 38.37 | 40.71 ($U_1=37.86$) |
| Peer 2 | 31.74 | 34.75 | 30.49 ($U_2=28.39$) |
| Peer 3 | 28.92 | 31.37 | 27.73 ($U_3=18.93$) |
| Free Rider | 33.95 | - | - |

a simple P2P system that cannot detect and cannot prevent free-riders (e.g., Gnutella [26]), (ii) a more sophisticated and adaptive P2P system that can detect and prevent free-riders effectively but does not consider the multimedia utility impact (e.g., a BitTorrent-like system [11]), and (iii) the proposed P2P system that explicitly considers each peer's multimedia attributes as well as free-rider prevention when performing resource reciprocation. Note that we present the result for a representative coalition that was formed in a network with 1000 peers.

We assume that the peers are interested in the *Stefan* video sequence. The maximum upload bandwidth of each peer is given by $L_i^{MAX}[\text{kbps}](1 \leq i \leq 3) = \{800, 600, 400\}$ and the unit costs of uploading rate are peer-dependent and pre-determined. In scenarios (i) and (ii), we assume that each peer shares part of its upload bandwidth with the matched peers due to the incurred cost. However, in scenario (iii), the upload bandwidth allocation is negotiated by PBS considering the utility impact. To study the impact of the resource reciprocation strategies on the utilities, we consider two simple

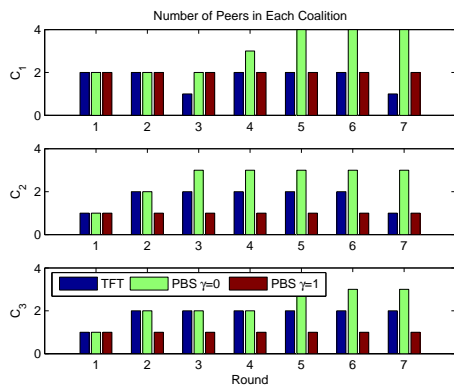
bandwidth allocation strategies. The available bandwidth can be equally divided (Equal Bandwidth Division) for scenarios (i) and (ii) as in a BitTorrent-like system, and equivalently $\mathbf{p}_E = (1/3, 1/3, 1/3)$ for scenario (iii). Alternatively, the available bandwidth can be proportionally divided (Proportional Bandwidth Division) based on each peer's maximum upload bandwidth for scenarios (i) and (ii), and thus, $\mathbf{p}_P = (4/9, 3/9, 2/9)$ for scenario (iii). We assume that a free-rider is connected to peer 1. Simulation results are shown in Table II.

The results of scenario (i) and scenario (ii) using the two bandwidth division strategies (equal bandwidth division and proportional bandwidth division) show that the quality derived in scenario (ii) is always higher than that derived in scenario (i). This occurs because in scenario (i) the free-rider obtains a part of the available bandwidth without contribution to the coalition. The results of scenario (ii) and scenario (iii) in the bandwidth division strategies show that scenario (iii) provides improved qualities since the PBS explicitly considers the impact of peers' attributes as well as the incurred cost. Moreover, scenario (iii) always provides improved individual qualities except for peer 2 and peer 3 in the proportional bandwidth division. The results show that the proportional bandwidth division strategy in scenario (i) and (ii) does not result in proportional utilities for the peers in the coalition. However, PBS does provide rewards, in terms of the utility, proportional to their contributions. These proportional rewards act as an incentive for the peers.

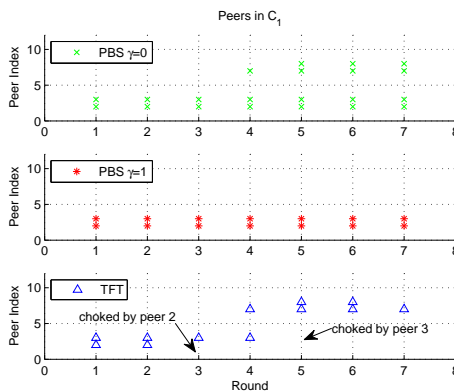
C. Comparison with TFT Strategy

In this section, we investigate the coalition formation and the resulting performance in terms of utilities based on the proposed resource reciprocation strategies (i.e., PBS with $\gamma = 0$ and $\gamma = 1$) and a TFT strategy deployed in P2P systems such as BitTorrent [11]. To compare both strategies, we assume that there is no cost for providing upload bandwidth in coalition based bandwidth sharing strategy, as the TFT strategy does not consider it. Moreover, since the TFT strategy focuses only on the downloading rates from other peers, the utility in the proposed solution is defined to represent the downloading rate. The inter-peer utility comparison ratio vector is determined based on the download rates. We also assume that the maximum number of parallel uploads for a peer deploying the TFT strategy is two, as peers have the maximum number of parallel uploads in the leecher state [11]. In the following simulations, a peer creates its coalition by sequentially joining 11 matched peers, and bandwidth negotiation is performed in each round (i.e., when peers join or leave the coalition). The peer is downloading the *Foreman* or *Mobile* video sequences and its maximum upload bandwidth is 850 Kbps, which is the highest upload bandwidth among coalition peers. Note that all peers in the coalition provide their maximum upload bandwidth, which is assumed to be constant in this experiment, as no cost is incurred for providing their upload bandwidth.

The simulation results are shown in Fig. 5, Fig. 6, and Table III. Fig. 5 shows the number and index of associated peers in selected coalitions (coalition 1, 2, and 3) over time (rounds). The corresponding achieved quality of peer 1 are



(a) Number of peers in each coalition



(b) Peer index in coalition 1

Fig. 5. Peers in each coalition for different strategies over time (rounds).

TABLE III
THE ACHIEVED QUALITIES WITH DIFFERENT STRATEGIES

| <i>Foreman</i> | Avg. PSNR | Max. PSNR | Min. PSNR | Difference |
|----------------------|-----------|-----------|-----------|------------|
| PBS ($\gamma = 0$) | 38.1 dB | 41.4 dB | 33.7 dB | 7.7 dB |
| PBS ($\gamma = 1$) | 39.2 dB | 41.1 dB | 36.5 dB | 4.6 dB |
| TFT | 35.2 dB | 41.1 dB | 28.0 dB | 13.1 dB |
| <i>Mobile</i> | Avg. PSNR | Max. PSNR | Min. PSNR | Difference |
| PBS ($\gamma = 0$) | 29.9 dB | 32.6 dB | 26.5 dB | 6.1 dB |
| PBS ($\gamma = 1$) | 31.4 dB | 32.3 dB | 29.5 dB | 2.8 dB |
| TFT | 26.6 dB | 32.1 dB | 21.6 dB | 10.5 dB |

shown in Fig. 6. During the first 6 seconds (the first round), the achieved quality of peer 1 are the same for all the strategies (i.e., TFT, PBS with $\gamma = 0$, and PBS with $\gamma = 1$). This occurs due to the small number of coalition peers (in this example, two peers – peer 2 and peer 3 – are associated) and that they provide their maximum upload bandwidth to peer 1. However, as more peers join this P2P network, and they create their own coalitions, the download rates of peer 1 vary. Since the TFT strategy allows peers to maintain a fixed maximum number of parallel uploads with peers providing the highest upload rates, its download rates can be smaller than the proposed coalition based resource reciprocation strategy (i.e., PBS with $\gamma = 0$). Moreover, due to the fixed number of parallel uploads in the TFT strategy, other peers can choke this peer, leading to significant loss of download rate and quality (e.g., the peer is choked by peer 2 and 3 at 20 second and at 50 second (round

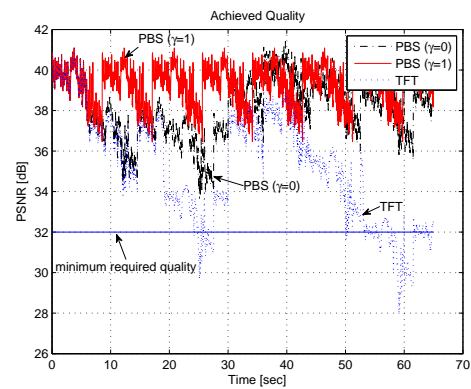
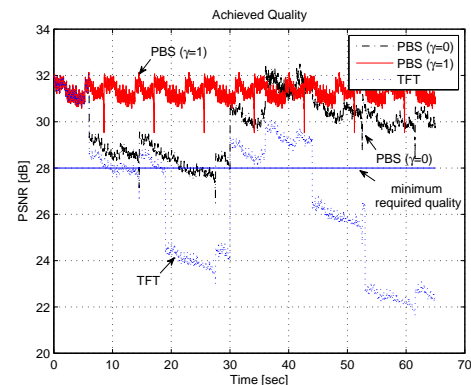
(a) *Foreman* sequence(b) *Mobile* sequence

Fig. 6. The achieved qualities for different strategies.

3 and 5). However, PBS does not abruptly deny to provide upload rates. Rather, it enables peers to scale the upload rates smoothly, leading to smooth changes of quality even in the case of the PBS with $\gamma = 0$.

We can also observe a period of time where peer 1's derived qualities decrease for both the TFT strategy and PBS with $\gamma = 0$ even though the same peers are associated (e.g., from 35 second to 53 second (at round 5, 6, and 7)). In this period, the associated peers for peer 1 create their own coalitions and negotiate their upload rates with their coalition peers. Hence, they decrease the upload rates for peer 1, resulting in decrease of download rates of peer 1. In the case of $\gamma = 1$, as we shown in the previous section, the quality can be derived at a similar level over the time. Note that PBS with $\gamma = 1$ ensures a minimum multimedia quality (in these simulations, 32dB PSNR for *Foreman* and 28dB PSNR for *Mobile* are assumed.), while the other strategies do not. Note that in Fig. 6, the upload bandwidth from each peer is fixed, but different PSNRs are achieved due to different sequence characteristics (e.g., *Foreman* and *Mobile*). However, if larger upload bandwidth is available for each peer, a higher PSNR can be derived.

These simulation results also show the resulting delay. If the minimum required quality is not satisfied, the peers may pause their displays until they find additional peers that have the necessary chunks. This can incur undesirable delay during the broadcasting besides the initial delay for the resource reciprocation process. In the simulation results, the TFT strategy

can incur significant delays during the broadcasting. Therefore, although the PBS incurs some initial delays for the resource reciprocation process, the overall delay can be reduced. This forms an interesting area for future research.

D. Application to Other Usage Scenarios

Simulations results presented in this section show that the proposed resource reciprocation based on the PBS with $\gamma = 1$ is preferred for a multimedia broadcasting system, which requires a stable multimedia quality. Hence, for example, the proposed approach can be adopted into proxy-based solutions (e.g., [31]) for several streaming video applications such as distance learning, Internet TV broadcasting, and video-on-demand (VoD) streaming services. Unlike P2P systems, where the service level can be determined by individual peers, if the proposed approach is implemented in the proxy servers, $\gamma = 1$ for PBS can be ensured, as they are dedicated to video distribution to clients. Thus, the proposed approach enables the proxy servers to support a stable quality multimedia services, while considering several constraints such as the different available bandwidths of heterogeneous clients.

The proposed approach can also be deployed in chaining schemes (e.g., [32]), as each client can forward its cached video streams to its associated clients. Hence, each client can divide its available bandwidth, while considering the quality impact as well as the associated clients' video characteristics.

VII. CONCLUSIONS

In this paper, we propose quality driven bandwidth reciprocation strategies for multimedia broadcasting over P2P networks. The proposed framework enables peers to create/join coalitions with which they can cooperatively share their available chunks. Moreover, the proposed approach based on PBS enables the peers to reciprocate their resources such that the resulting utilities are proportional to their contributions to the coalition value. In order to investigate the impact of the peers' attitude toward the other peers on the derived multimedia quality, the service level is introduced. We show that peers with a higher service level can guarantee more stable multimedia quality in the coalition. This feature is important for multimedia broadcasting applications. Simulation results show that the proposed coalition-based resource reciprocation strategy is more efficient for P2P multimedia broadcast than the TFT strategy currently deployed in BitTorrent-like networks. We also discuss how the proposed approach can be implemented in other video streaming applications.

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